Record List Display

INVENTOR-INFORMATION:

NAME CITY STATE COUNTRY RULE-47

Kunitake, SetsuNakai-machiJPYamashita, IchiroNakai-machiJPKawabe, ShigehisaNakai-machiJP

US-CL-CURRENT: 707/517; 707/514, 707/531

Full Title Citation Front Review Classification Date Reference Sequences Attachments Claims RMC Draw Desc Image

9. Document ID: US 20010007987 A1

L9: Entry 9 of 11 File: PGPB

Jul 12, 2001

PGPUB-DOCUMENT-NUMBER: 20010007987 PGPUB-FILING-TYPE: new-utility

DOCUMENT-IDENTIFIER: US 20010007987 A1

TITLE: Structured-document search apparatus and method, recording medium storing structured-document searching program, and method of creating indexes for searching

structured documents

PUBLICATION-DATE: July 12, 2001

INVENTOR - INFORMATION:

NAME

CITY STATE COUNTRY RULE-47

Igata, Nobuyuki Kawasaki JP

US-CL-CURRENT: 707/3; 707/1

Full Title Citation Front Review Classification Date Reference Sequences Attachments Claims KWC Draw Desc Image

10. Document ID: US 6424979 B1

L9: Entry 10 of 11 File: USPT

Jul 23, 2002

US-PAT-NO: 6424979

DOCUMENT-IDENTIFIER: US 6424979 B1

TITLE: System for presenting and managing enterprise architectures

DATE-ISSUED: July 23, 2002

INVENTOR-INFORMATION:

NAME CITY STATE ZIP CODE COUNTRY Livingston; Eric Bruce Fairfax VA

Vogt; Lisa Conners Bethesda MD Gantz; Stephen David Arlington VA Zipkin; David Samuel Arlington VA Richards; Matthew Todd Vienna VA Rocchio; Sheila Caroline Arlington VA

ASSIGNEE-INFORMATION:

NAME CITY STATE ZIP CODE COUNTRY TYPE CODE

American Management Systems, Inc. Fairfax VA 02



DATE FILED: December 30, 1998

PARENT-CASE:

CROSS REFERENCE TO RELATED APPLICATION This application is related to U.S. application entitled CONTENT MANAGEMENT SYSTEM having Ser. No. 09/222,831 by Baxter and Voqt, filed Dec. 30, 1998, now U.S. Pat. No. 6,356,903 and incorporated by reference herein.

INT-CL: [07] G06 F 17/00

US-CL-ISSUED: 707/511; 707/203, 707/501.1 US-CL-CURRENT: 707/511; 707/203, 707/501.1

FIELD-OF-SEARCH: 707/10, 707/1, 707/2, 707/3, 707/100, 707/101, 707/104, 707/9, 707/104.1, 707/501.1, 707/203, 707/511, 709/202, 714/100

PRIOR-ART-DISCLOSED:

U.S. PATENT DOCUMENTS

PAT-NO	ISSUE-DATE	PATENTEE-NAME	US-CL
5263129	November 1993	Ikegaya et al.	
5273434	December 1993	Peck	
5311424	May 1994	Mukherjee et al.	705/29
5317729	May 1994	Mukherjee et al.	707/3
5455407	October 1995	Rosen	235/380
<u>5737739</u>	April 1998	Shirley et al.	
5842196	November 1998	Agarwal et al.	707/2
5862346	January 1999	Kley et al.	709/200
<u>5907837</u>	May 1999	Ferrel et al.	707/3
<u>5949876</u>	September 1999	Ginter et al.	380/4
5970476	October 1999	Fahey	705/28
6067559	May 2000	Allard et al.	709/202
6078924	June 2000	Ainsbury et al.	707/101
6085220	July 2000	Courts et al.	709/201
6119229	September 2000	Martinez et al.	713/200
6151601	November 2000	Papierniak et al.	707/10
6157915	December 2000	Bhaskaran et al.	705/ 7
6182091	January 2001	Pitkow et al.	707/501

OTHER PUBLICATIONS

E. James Whitehead, Jr. "World Wide Web distributed authoring and versioning (WebDAV): an introduction", StandardView vol. 5, No. 1, Mar. 1997, pp. 3-8.*

U.S. patent application Ser. No. 09/222,831, Baxter et al., filed Dec. 30, 1998.

Extensible Markup Language (XML) 1.0, W3C Recommendation Feb. 10, 1998.

Texcel International A.B., Texcel Information Manager Release 2.0 Product Brief, 1998. Delphi Consulting Group, Delphi Opinion, Astoria 3.0 Chrystal Software, Jun. 1998.

Arbortext, Inc., Introducting Epic.TM. An Enterprise Solution for the Product Information Chain, 1998.

Microstar Software Inc., Content Management Systems.

AIS Software, Application Sheet, Merging XML Content with Relational Databases On-line Catalog Publishing with Dual Prism.

AIS Software, News Release, DUAL PRISM.TM. Now Shipping!, Official Launch of Dual Prism in North America at XML Chicago 1998, Oct. 1998.

AIS Software, Dual Prism.RTM. Powerful Web Content Management and Online Delivery.

ART-UNIT: 2171

PRIMARY-EXAMINER: Metjahic; Safet

ASSISTANT-EXAMINER: Le; Uyen

ŧ

ABSTRACT:

A system providing views of a technical architecture of an enterprise that take into account the content type interest, level of detail and time frame of desired information. A user, with a browser interface, selects a portal specifying the type of content to be viewed. A profile of the user indicates the level of detail and time frame of information to be obtained. The interface provides the user's selection of desired information within the portal in the form of a page request that is converted into queries of a database that seek content satisfying the type, level of detail and time frame attributes of the request. The content is formed into a web page and provided to the interface by the web server. Content is separated into atomic units allowing the information to be rearranged responsive to each user's needs.

28 Claims, 26 Drawing figures

Full Title Citation Front Review Classification Date Reference Sequences Attachments Claims RMC Draw Desc Image

STATE

STATE

11. Document ID: US 6263332 B1

L9: Entry 11 of 11

File: USPT

ZIP CODE

ZIP CODE

Jul 17, 2001

US-PAT-NO: 6263332

DOCUMENT-IDENTIFIER: US 6263332 B1

TITLE: System and method for query processing of structured documents

DATE-ISSUED: July 17, 2001

INVENTOR-INFORMATION:

NAME

CITY Nasr; Roger I.

Austin ΥY Austin

Webber; Neil

ASSIGNEE-INFORMATION:

NAME

CITY

Vignette Corporation Austin

ТX

COUNTRY TYPE CODE

02

COUNTRY

APPL-NO: 09/ 134263

DATE FILED: August 14, 1998

INT-CL: [07] G06 F 17/30

US-CL-ISSUED: 707/5; 707/104, 707/513 US-CL-CURRENT: <u>707/5</u>; <u>707/104.1</u>, <u>707/513</u>

FIELD-OF-SEARCH: 707/1-4.5, 707/513, 707/523

PRIOR-ART-DISCLOSED:

U.S. PATENT DOCUMENTS

PAT-NO	ISSUE-DATE	PATENTEE-NAME	US-CL
5138615	August 1992	Lamport et al.	370/400
5649186	July 1997	Ferguson	707/10
5655130	August 1997	Dodge et al.	707/511
5752021	May 1998	Nakatsuyama et al.	707/5
5778400	July 1998	Tateno	707/513
5875441	February 1999	Nakatsuyama	707/1
5893109	April 1999	DeRose et al.	707/104
5920879	July 1999	Kyojima et al.	707/517
5940822	August 1999	Haderle et al.	707/3
5940842	August 1999	Sakuta	707/513
5956726	September 1999	Aoyama et al.	707/102
6009436	December 1999	Motoyama et al.	707/102
6012098	January 2000	Bayeh et al.	709/246
6018710	January 2000	Wynblatt et al.	704/260
6023714	February 2000	Hill et al.	707/513
6047280	April 2000	Ashby et al.	707/2
6055538	April 2000	Kessenich et al.	707/101
6072801	June 2000	Wood, Jr. et al.	370/437
6098071	August 2000	Aoyama et al.	707/102
6108676	August 2000	Nakatsuyama	707/522
6128617	October 2000	Lowry	707/100

FOREIGN PATENT DOCUMENTS

FOREIGN-PAT-NO	PUBN-DATE	COUNTRY	US-CL
0 601 550 A2	July 1993	EP	

OTHER PUBLICATIONS

Marden, Philip, and Munson, Ethan. "Multiple Presentations of WWW Documents Using Style Sheets", Proceedings of the Workshop on New Paradigms in Information Visualization and Manipulation, Nov. 10-14, 1997, pp. 75-78.*

"Microsoft Press Computer Dictionary Third Edition", Definition of the word `tree`, 1997, p. 477.*

Lee, Y.K., Yoo, S-J, and Yoon, K. "Index Structures for Structured Documents", Proceedings of the First ACM International Conference on Digital Libraries, Mar. 20-23, 1996, pp. 91-99.*

Aho, A.V., Hopcraft, J.E., and Ullman, J.D. "Data Structures and Algorithms", Reading, Addison Wesley, 1983. pp. 78-82. QA76.9.D35A38.*

Thompson, Henry S.: "An Introduction to XSL: Slides from a talk to SGML UK" from http://www.itg.ed.ac.uk/.about.ht/swindon.html. Oct. 27, 1997. Sections 5,6,9.*
Lee, Young Kyu: "Querying Structured Documents". System Sciences, 1996., Proceedings of the Twenty-Ninth Annual Hawaii Conference on System Sciences, vol. 2, pp. 155-164.

ART-UNIT: 217

PRIMARY-EXAMINER: Homere; Jean R.

ASSISTANT-EXAMINER: Wassum; Luke S

ABSTRACT:

A computer-implemented method of retrieving information in a first markup language through a query engine and presenting the information in any required markup language. A user inputs a query and may invoke a number of transformative sequences. These sequences contain a markup language pattern and an action, which may include transforming the tags in the first markup language to tags in a different markup language. The appropriate transformative sequence is selected and the pattern from the transformative sequence is compiled. The compiled pattern is used to perform rapid and



efficient searches of documents in the database. A predicate check using the binary coding of the node as well as ancestor information confirms the node. The leaf information associated with a confirmed node is then stored. If necessary, the action from the transformative sequence is applied to change the markup language of the leaf information to that of the user.

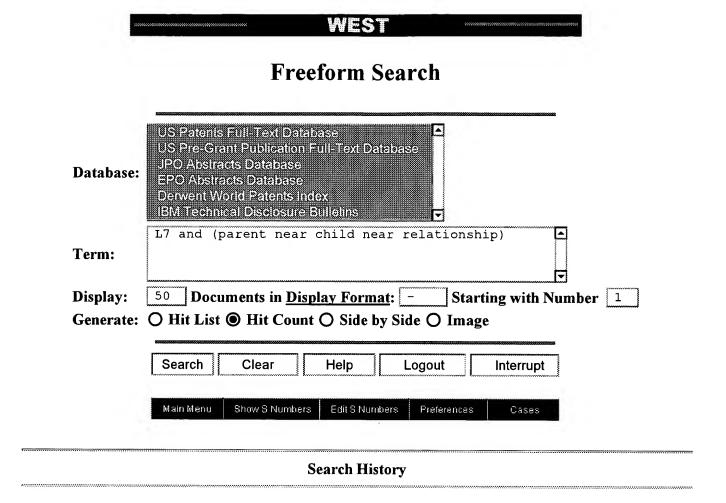
36 Claims, 10 Drawing figures

Full Title Citation Fro	Review Classification Date Reference Sequences Affachments Claims MMC Draw De	sc Image
	Generate Collection Print	

Term	Documents
PARENT.DWPI,TDBD,EPAB,JPAB,USPT,PGPB.	173200
PARENTS.DWPI,TDBD,EPAB,JPAB,USPT,PGPB.	11498
CHILD.DWPI,TDBD,EPAB,JPAB,USPT,PGPB.	116384
CHILDREN.DWPI,TDBD,EPAB,JPAB,USPT,PGPB.	68795
CHILDRENS.DWPI,TDBD,EPAB,JPAB,USPT,PGPB.	2859
CHILDS.DWPI,TDBD,EPAB,JPAB,USPT,PGPB.	4351
RELATIONSHIP.DWPI,TDBD,EPAB,JPAB,USPT,PGPB.	838370
RELATIONSHIPS.DWPI,TDBD,EPAB,JPAB,USPT,PGPB.	133990
(7 AND (PARENT NEAR CHILD NEAR RELATIONSHIP)).USPT,PGPB,JPAB,EPAB,DWPI,TDBD.	11
(L7 AND (PARENT NEAR CHILD NEAR RELATIONSHIP)).USPT,PGPB,JPAB,EPAB,DWPI,TDBD.	11

Display Format: - Change Format

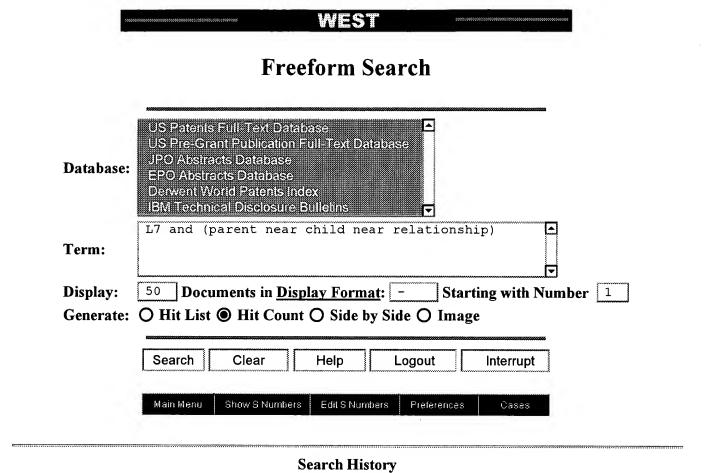
Previous Page Next Page



DATE: Friday, July 26, 2002 Printable Copy Create Case

et Name Query			Set Name
de by sid	e		result set
DB=U	SPT,PGPB,JPAB,EPAB,DWPI,TDBD; PLUR=YES; OP=OR		
<u>L9</u>	L7 and (parent near child near relationship)	11	<u>L9</u>
<u>L8</u>	L7 and (flat near data near structur\$)	1	<u>L8</u>
<u>L7</u>	L6 and (data near structur\$)	316	<u>L7</u>
<u>L6</u>	((707/\$)!.CCLS.) and xml	698	<u>L6</u>
<u>L5</u>	L1 and xml	0	<u>L5</u>
<u>L4</u>	L1 and (flat near data near structur\$)	0	<u>L4</u>
<u>L3</u>	L2 and (flat)	0	<u>L3</u>
<u>L2</u>	L1 and (convert\$ and (hierarch\$ near data near structur\$))	11	<u>L2</u>
<u>L1</u>	(4954969 5454101 5856925 5924087 6055515 4521882 4928247 4995112 5224100 5303367 5392220 5430870 5438511 5842213 5923881 5960433 6112209 4555771 4796199 4864497 4947447 5251322 5265244 5280547 5295261 5341466 5363320 5394249 5410675 5467471 5490258 5495440 5506788 5542085 5561608 5577239 5627999 5634123 5636158 5648813 5651101 5664177 5668965 5701467 5703800 5761637 5771384 5784526 5799183 5808901).pn.	135	<u>L1</u>

END OF SEARCH HISTORY



DATE: Friday, July 26, 2002 Printable Copy Create Case

et Name	<u>e Query</u>	Hit Count	<u>Set Name</u>
de by sid	e		result set
DB=U	SPT,PGPB,JPAB,EPAB,DWPI,TDBD; PLUR=YES; OP=OR		
<u>L9</u>	L7 and (parent near child near relationship)	11	<u>L9</u>
<u>L8</u>	L7 and (flat near data near structur\$)	1	<u>L8</u>
<u>L7</u>	L6 and (data near structur\$)	316	<u>L7</u>
<u>L6</u>	((707/\$)!.CCLS.) and xml	698	<u>L6</u>
<u>L5</u>	L1 and xml	0	<u>L5</u>
<u>L4</u>	L1 and (flat near data near structur\$)	0	<u>L4</u>
<u>L3</u>	L2 and (flat)	0	<u>L3</u>
<u>L2</u>	L1 and (convert\$ and (hierarch\$ near data near structur\$))	11	<u>L2</u>
<u>L1</u>	(4954969 5454101 5856925 5924087 6055515 4521882 4928247 4995112 5224100 5303367 5392220 5430870 5438511 5842213 5923881 5960433 6112209 4555771 4796199 4864497 4947447 5251322 5265244 5280547 5295261 5341466 5363320 5394249 5410675 5467471 5490258 5495440 5506788 5542085 5561608 5577239 5627999 5634123 5636158 5648813 5651101 5664177 5668965 5701467 5703800 5761637 5771384 5784526 5799183 5808901).pn.	135	<u>L1</u>

END OF SEARCH HISTORY

End of Result Set

Generate Collection Print

L2: Entry 11 of 11

File: DWPI

Mar 15, 1994

DERWENT-ACC-NO: 1994-091836

DERWENT-WEEK: 199411

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TITLE: Data record to hybrid structure conversion method for database - storing navigational data in topological tree structure with each leaf node specifying unique record in relational informational database

Patent Serial Number (1): 5295261

Basic Abstract Text (1):
The method involves separating the number of fields into navigational fields and informational fields, where geographic values encompassed in each navigational field of the data records are of a geographical class defined by the field type of the navigational field. The navigationa l fields are ordered into a number of distinct hierarchical data structures each comprising a number of node classes and links. Among the number of distinct hierarchical data structures one of the distinct hierarchical data structures is identified having at least one geographica 1 class in common.

Portions of the distinct hierarchical data structures having one or more geographic classes in common are merged by linking the structures to a single node class.

Basic Abstract Text (2):

For each the node class in the merged hierarchical data structure a table is created comprising an identifier describing the geographical class represented by the node class, an enumeration list encompassing the geographical values of the node class, and a number of linkage structures.

Standard Title Terms (1):

DATA RECORD HYBRID STRUCTURE CONVERT METHOD DATABASE STORAGE NAVIGATION DATA TOPOLOGICAL TREE STRUCTURE LEAF NODE SPECIFIED UNIQUE RECORD RELATED DATABASE

WEST

Generate Collection

Print

Search Results - Record(s) 1 through 11 of 11 returned.

Document ID: US 20020091702 A1

L9: Entry 1 of 11

File: PGPB

Jul 11, 2002

PGPUB-DOCUMENT-NUMBER: 20020091702

PGPUB-FILING-TYPE: new

DOCUMENT-IDENTIFIER: US 20020091702 A1

TITLE: Dynamic object-driven database manipulation and mapping system

PUBLICATION-DATE: July 11, 2002

INVENTOR-INFORMATION:

NAME

CITY

STATE

COUNTRY

RULE-47

Mullins, Ward

San Francisco

CA

US

US-CL-CURRENT: 707/100

Full Title Citation Front Review Classification Date Reference Sequences Attachments Claims KMC Draw Desc Image

2. Document ID: US 20020083034 A1

L9: Entry 2 of 11

File: PGPB

Jun 27, 2002

PGPUB-DOCUMENT-NUMBER: 20020083034

PGPUB-FILING-TYPE: new

DOCUMENT-IDENTIFIER: US 20020083034 A1

TITLE: Method and apparatus for extracting data objects and locating them in virtual

space

PUBLICATION-DATE: June 27, 2002

INVENTOR - INFORMATION:

NAME

CITY

STATE

COUNTRY

RULE-47

Orbanes, Julian Guzman, Adriana

Cambridge Cambridge MA MA US US

US-CL-CURRENT: 707/1

Full Title Citation Front Review Classification Date Reference Sequences Attachments Claims KMC Draw Desc Image

3. Document ID: US 20020078054 A1

L9: Entry 3 of 11

File: PGPB

Jun 20, 2002

PGPUB-DOCUMENT-NUMBER: 20020078054

PGPUB-FILING-TYPE: new

DOCUMENT-IDENTIFIER: US 20020078054 A1

TITLE: Group forming system, group forming apparatus, group forming method, program,

and medium

PUBLICATION-DATE: June 20, 2002

INVENTOR-INFORMATION:

NAME

CITY Osaka STATE

COUNTRY

RULE-47

Kudo, Takahiro

Kyotanabe-shi

JР

Matsuura, Satoshi

ny o canabe

JР

Ozawa, Jun

Nara-shi

JΡ

US-CL-CURRENT: 707/10

Full Title Citation Front Review Classification Date Reference Sequences Attachments Claims KMC Draw Desc Image

4. Document ID: US 20020069215 A1

L9: Entry 4 of 11

File: PGPB

Jun 6, 2002

PGPUB-DOCUMENT-NUMBER: 20020069215

PGPUB-FILING-TYPE: new

DOCUMENT-IDENTIFIER: US 20020069215 A1

TITLE: Apparatus for viewing information in virtual space using multiple templates

PUBLICATION-DATE: June 6, 2002

INVENTOR-INFORMATION:

NAME

CITY

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COUNTRY

RULE-47

Orbanes, Julian

Cambridge

MA

US

Guzman, Adriana

Cambridge

MA

US

US-CL-CURRENT: <u>707</u>/<u>500</u>

Full | Title | Citation | Front | Review | Classification | Date | Reference | Sequences | Attachments | Claims | KWIC | Draw Desc | Image

5. Document ID: US 20020035555 A1

L9: Entry 5 of 11

File: PGPB

Mar 21, 2002

PGPUB-DOCUMENT-NUMBER: 20020035555

PGPUB-FILING-TYPE: new

DOCUMENT-IDENTIFIER: US 20020035555 A1

TITLE: System and method for building and maintaining a database

PUBLICATION-DATE: March 21, 2002

INVENTOR-INFORMATION:

NAME

CITY

STATE

COUNTRY

RULE-47

Wheeler, David B.

Austin

TX

US

Wotring, Steven C.

Austin

TX

US

US-CL-CURRENT: 707/1; 707/104.1

Full | Title | Citation | Front | Review | Classification | Date | Reference | Sequences | Attachments | Claims | MMC | Draw Desc | Incape

6. Document ID: US 20010049688 A1

L9: Entry 6 of 11

File: PGPB

Dec 6, 2001

PGPUB-DOCUMENT-NUMBER: 20010049688

PGPUB-FILING-TYPE: new

DOCUMENT-IDENTIFIER: US 20010049688 A1

TITLE: System and method for providing an intelligent multi-step dialog with a user

PUBLICATION-DATE: December 6, 2001

INVENTOR-INFORMATION:

NAME	CITY	STATE	COUNTRY	RULE-47
Fratkina, Raya	Yorktown Heights	NY	US	
Anderson, Monica	San Jose	CA	US	
Angel, Mark A.	Napa	CA	US	
Copperman, Max	Santa Cruz	CA	US	
Huffman, Scott B.	Redwood City	CA	US	
Kay, David	Los Gatos	CA	US	
Stern, Robert	Cupertino	CA	US	

US-CL-CURRENT: 707/104.1; 345/968

Full Title Citation Front Revi	w Classification Date Referenc	⊵ Sequences Attachments Claim	is FWMC Draw Desc Image

7. Document ID: US 20010037345 A1

L9: Entry 7 of 11

File: PGPB

Nov 1, 2001

PGPUB-DOCUMENT-NUMBER: 20010037345

PGPUB-FILING-TYPE: new

DOCUMENT-IDENTIFIER: US 20010037345 A1

TITLE: Tagging XML query results over relational DBMSs

PUBLICATION-DATE: November 1, 2001

INVENTOR-INFORMATION:

NAME CITY STATE COUNTRY RULE-47 Kiernan, Gerald George US San Jose CA Shanmugasundaram, Jayavel San Jose CA US Shekita, Eugene Jon San Jose CA US

US-CL-CURRENT: <u>707/513</u>

ron i me Citation Front Neview Classification	on i Date i Kererence i Sequences i Attachments i Chaims i	KWWC Draw Desc Image

S. Document ID: US 20010	0018697 A1	
L9: Entry 8 of 11	File: PGPB	Aug 30, 2001

PGPUB-DOCUMENT-NUMBER: 20010018697

PGPUB-FILING-TYPE: new

DOCUMENT-IDENTIFIER: US 20010018697 A1

TITLE: Structured document processing system and structured document processing method

PUBLICATION-DATE: August 30, 2001

Generate Collection

Print

L2: Entry 10 of 11

File: USPT

DOCUMENT-IDENTIFIER: US 4555771 A

TITLE: Data processing system for data base management of matrix type data

<u>US PATENT NO.</u> (1):

Brief Summary Text (9):

Another object of the invention is to provide a data processing system for virtually converting matrix type data to a desired hierarchical structure data.

Brief Summary Text (13):

wherein al, a2 . . . an is the n-dimensional coordinate and A1, A2 . . . An is the coordinate axis size. The calculator includes a converter connected to the scheme code storage, the address storages and the designating device for providing correspondence between the contents of the designated one of the address storages and the n-dimensional coordinate in accordance with the contents of the scheme code storage, the converter having an output, and a linear expander connected to the n-coordinate axis size storage, the designating device and the converter for executing operation of the one-dimensional address information value in accordance with the output of the converter and the contents of the n-coordinate axis size storage.

Drawing Description Text (14):

FIG. 8F is a block diagram of an embodiment of a converter unit CV;

<u>Detailed Description Text (43):</u>

FIG. 8F is a block diagram of an embodiment of a <u>converter</u> unit CV, which distributes signal inputs supplied to input terminals I1, I2 and I3 to output terminals O1, O2 and O3 in accordance with an input supplied to terminal C when a pulse is provided at terminal E. More particularly, the input and output terminals of the <u>converter</u> unit are connected, for example, I1 to O1, I2 to O2 and I3 to O3 when a specific value is applied to the terminal C, and are connected, for example, I1 to O2, I2 to O3 and and I3 to O1 when a different value is applied to the terminal C. When a terminal D is driven, an inverted replacement is provided. The <u>converter</u> unit CV is operated with a delay of one cycle from the driving pulse. When "O" is applied to any of the input terminals I1, I2 and I3, a terminal N is driven. I1 to I3, O1 to O3 and C are multi-bit signals and E and N are single-bit signals.

Detailed Description Text (44):

The converter unit CV of FIG. 8F includes a decoder 601, AND gates 602 to 610, OR gates 611 to 615, latches 616 to 618, a delay circuit 619 for a one cycle delay and zero detect circuits 620 to 622. When the input C equals zero, for example, the decoder 601 drives its output line L1, and accordingly the input terminals I1, I2 and I3 are connected to the OR gates 611, 612 and 613 via the AND gates 602, 605 and 608, respectively.

<u>Detailed Description Text</u> (52):

In FIG. 9, 101 to 119 are register units R, as shown in FIG. 8C, 120 to 122 are arithmetic units AU as shown in FIG. 8A, 123 to 128 are bypass gates BG, as shown in FIG. 8B, 129 to 131 are converter units CV, as shown in FIG. 8F, 132 is a set selection unit SS, as shown in FIG. 8D, 133 and 134 are OR gates OR, as shown in FIG. 8E, 136 and 137 are NOT circuits NOT, as shown in FIG. 8H, 135 and 138 are AND gates AND, as shown in FIG. 8I, 139 is a linear expander LE, as shown in FIG. 8G and 140 is a compose matrix unit CM, as shown in FIG. 8J.

Detailed Description Text (53):

The registers 101, 102 and 103 store the sizes A1, A2 and A3, respectively, of the

three coordinates a1, a2 and a3, respectively. In FIG. 1, specifically, A1=2, A2=2, A3=3. The register 119 stores the scheme code C, which designates the desired hierarchical structure. The converter 129 replaces the size information A1 to A3 received from the registers 101 to 103, respectively, and distributes such information to the registers 104 to 106, respectively, in accordance with the code C provided by the register 119. When the scheme in FIG. 4 is designated, for example, A2 is stored in the register 104, A3 is stored in the register 105 and A1 is stored in the register

Detailed Description Text (58):

The code for designating the third set SS3 is then fed to the set selection unit 132 as the code S in order to operate the arithmetic unit 122. The contents of the register 109 are thus incremented by the modulus on the basis of the contents of the register 106. When the process is completed, a pulse is provided at the terminal K of the arithmetic unit 122, and is fed to the terminal E of the converter 130 and to the terminal Q of the linear expander 139. The converter 130 transfers the signals I1, I2 and I3 received from the registers 109, 111 and 112 to its outputs 01, 02 and 03, in accordance with the scheme designation code C received from the register 119. The linear expander 139 provides the aforedescribed process. In the current state, access is to the #3 record (113) in FIG. 4, for example. In other words, I1=1, I2=1 and I3=3 as the inputs of the linear expander 139, and A1=2, A2=2 and A3=3. Therefore, ##EQU2## More particularly, the one-dimensional address #3 may be obtained from the three-dimensional coordinate (113).

Detailed Description Text (64):

It is possible to designate a specific member record of the other owner record in a known data base management system under the condition that a specific member record is designated under a specific owner record. The location information of the new owner record 0' may be set in the register 117. The one-dimensional address P' of a new member record may be set in the register 118. The contents of the register 118 are converted to a three-dimensional coordinate by the compose matrix unit 140. Each coordinate value of such three-dimensional coordinate is replaced by the converter 131 in accordance with the scheme designation code in the register 119 and set in each register 109, 111 and 112.

<u>Detailed Description Text</u> (66):

Furthermore, when the <u>converter</u> unit 130 detects that any of the inputs I1 to I3 is zero, a pulse is provided at its terminal N and fed to the terminal I2 of the AND gate 135. When the contents of any of the registers 109, 111 and 112 becomes zero, therefore, under the condition that a pulse is not output at the terminal E of the arithmetic unit 122, a signal N may be output at the terminal O of the AND gate 135. The signal N is an error signal.

CLAIMS:

calculating means including a converter means including a converter means, said converter means connected to said n-coordinate axes, said scheme code storage means, said address storage means and said selecting means for calculating the one-dimensional address from the contents of said address storage means, the contents of said scheme code storage means and the contents of said n-coordinate axes, said file means being connected to said calculating means and being accessed by the one-dimensional address output from said calculating means, said converter means providing correspondence between the contents of said selected one of said address storage means and said n-dimensional coordinate in accordance with the contents of said scheme code storage means said converter means including a decoder, responsive to input scheme code signals, for providing different outputs of said converter means for the same inputs from said n-coordinate-axes, said converter means providing an output, and linear expander means for executing operation of said one-dimensional address information value in accordance with the output of said converter means and the contents of said n-coordinate axes, said linear expander means connected to said n-coordinate axes, said selecting means, and said converter means.

L2: Entry 6 of 11 File: USPT

DOCUMENT-IDENTIFIER: US 5438511 A TITLE: Disjunctive unification

<u>US PATENT NO.</u> (1): 5438511

Brief Summary Text (3):

A number of general approaches have been proposed for solving a system such as a logical formula whose terms are equations. Some of these, such as term rewriting, operate directly on the terms, which are typically conjoined equations. Others, such as unification or merging, operate on terms converted into suitable data structures. Approaches of either of these types typically perform poorly if the system being solved includes disjunctions.

Brief Summary Text (18):

A closely related aspect of the invention is based on the recognition of a problem in associating context identifiers within a hierarchical data structure. In general, the unification of different context identifiers results in an expansion or cross product multiplication to consider all possible combinations of contexts, similar to expanding to the disjunctive normal form. Therefore, context identifiers that are relatively high in hierarchical data structures being unified result in an inefficient expansion between the context identifiers in the data structures.

Brief Summary Text (19):

This aspect of the invention solves this problem based on the recognition that the number of context identifiers that are relatively high in a hierarchical data structure can be reduced by combining context identifiers. Context identifiers can only be combined, however, if data units that depend from them can be combined. Therefore, this solution depends on including context identifiers in data structures such that the data units that depend from them can frequently be combined, either by unification or other techniques.

Detailed Description Text (17):

FIG. 1 shows a sequence of general steps, beginning with disjunctive system 10 that includes disjunctions 12 and 14, each illustratively including a number of feature-value equations although the invention is also applicable to other types of disjunctions, as described above. Each disjunction is converted to a data structure, and these data structures are then unified to obtain a solution of the disjunctive system.

Detailed Description Text (21):

The conversion of each disjunction into a data structure for unification could be done in various ways. For example, data structures 22 and 24 are <a href="https://nierarchical.com/hierarchical

Finally, the lowest level of each data structure includes feature-value pairs, each feature-value pair including a feature token and a value token. The feature-value pairs

Detailed Description Text (22):

As discussed above, unifying data structures 22 and 24 would be inefficient, because the context identifiers are relatively high in the hierarchy of each data structure, which would result in expansion to consider all possible combinations of context identifiers. Therefore, in accordance with a further aspect of the invention, alternative context-value pairs in data structures 22 and 24 are combined to produce data structures 32 and 34, in which each feature that previously occurred in a feature-value pair now has a disjunctive value data unit that includes appropriate context identifiers, so that the feature tokens are now above the context identifiers in the hierarchy.

within each disjunct depend from that disjunct's context identifier.

Detailed Description Text (40):

SOLVEAND handles each conjunct according to its relation, which may illustratively be equality or disjunction. The step in box 102 branches based on this relation. If the relationship is equality, SOLVEAND calls LOCATE to find the current values of the arguments and to convert equations into data structures; SOLVEAND then calls UNIFY to unify the data structures corresponding to the current values of the arguments, in box 104. In converting to data structures, LOCATE may in turn call LOCATE.VAR or LOCATE.AE, either of which may call CREATECONTEXTVALUE to create an entity of type Disjunction, which is a set of context-value pairs and therefore is a disjunctive value data unit as described above. If, rather than equality, the relationship is disjunction, SOLVEAND calls SOLVEOR to handle it, which begins an inner iterative loop in box 110. When unification in box 104 is complete or all the disjuncts have been handled in the inner iterative loop, SOLVEAND begins another iteration with the next conjunct, in box 100.

Detailed Description Text (52):

FIG. 6B shows f-structure 300 which is equivalent to disjunctive value 200, but within which each feature has a disjunctive value. The f-structure 300 may be thought of as a reduced form of disjunctive value 200. It could be obtained by converting disjunctive value 200 to its reduced form by combining the context-value pairs shown in FIG. 6A, a technique discussed in greater detail below; it could alternatively be obtained directly from disjunction (1). Like disjunctive value 200, it has outer brackets 302 and 304, but rather than having three terms within its outer brackets, f-structure 300 has two terms-one each for the two feature tokens A and B that occur in disjunction (1). In addition to a respective feature token, each of the terms within outer brackets 302 and 304 also includes a set of inner brackets 310 and 312. The inner bracket in each term contains a disjunctive value for the respective feature token.

Detailed Description Text (78):

The test in box 410 determines whether D1 and D2 are of type Disjunction. If either entity is not of type Disjunction, the step in box 412 calls COERCETODISJUNCTION to convert it into an entity of type Disjunction, which can be done by copying the non-disjunctive entity and then smashing the original to be an entity of type Disjunction with a single context-value pair, the context being the NULL context and the value being the copy of the non-disjunctive entity.

Detailed Description Text (83):

When all the pairs in D2 have been handled, UNIFY.CONTEXTS returns to the test in box 420 to handle the next context-value pair in D1. When all the pairs in D1 have been handled, the step in box 450 calls REDUCEDISJUNCTION, discussed in more detail below, to convert the unified version of D1 and D2 to the reduced form with disjunctive values. Then, the test in box 452 determines whether a swap was performed in box 416. If not, D1 is returned, in box 406, but if so, D2 is returned in box 454.

Detailed Description Text (135):

Although the invention does not solve the NP-complete problem, it does avoid exponential growth for typical examples of the problem that are decomposible into independent P class problems. For example, in a long sentence, different disjunctions are typically independent if far apart, so that use of the disjunctive normal form is unnecessary and the disjunctions can be treated independently as P class problems. In cases where disjunctions are not independent, the use of context identifiers keeps track of their relationships and allows the consideration of combinations of disjuncts from different disjunctions as necessary without solving the NP-complete problem. Since any NP-complete problem can be converted into any other, the invention should be useful with other problems in which disjunctions occur that can be treated as independent P

class problems.

L2: Entry 9 of 11 File: USPT

DOCUMENT-IDENTIFIER: US 5295261 A

TITLE: Hybrid database structure linking navigational fields having a hierarchial database structure to informational fields having a relational database structure

<u>US PATENT NO.</u> (1): 5295261

Drawing Description Text (3):

FIGS. 2(A), 2(B), and 2(C) illustrate the manner in which one of the tables shown in FIG. 1 can be converted into a database according to the present invention.

Detailed Description Text (17):

The final step of converting the database consists of replacing table 41 by a tree structure 50 as shown in FIG. 2(C). The nodes in tree structure 50 are divided into sets shown at 51-54. Each set of nodes corresponds to a column in table 41. The nodes in set 51 correspond to the state, those in set 52 correspond to the city, those in set 53 correspond to the street address, and those in set 54 correspond to the unique identifier defined for each record in the original database. Each node in a given set is linked to a node in a set one level higher up in the hierarchy. Hence, there is no need to repeat the redundant data. For example, by traversing the tree from any given city node to the state node to which it is linked, one may ascertain the state in which the city is located.

CLAIMS:

1. A method for <u>converting</u> organized data from a data record structure into a hybrid structure which allows for efficient storage and retrieval of the organized data by a digital computer, where the data record structure includes a plurality of records each divided into a plurality of fields, where each field of the plurality of fields stores data of a field type associated with the field and where the organized data in the data record structure is encompassed in the plurality of fields of the plurality of records, the method comprising the steps of:

ordering said navigational fields into a plurality of distinct hierarchical data structures each comprising a plurality of node classes and links therebetween, where each said node class corresponds to a specific geographical class and each node in said node class corresponds to a distinct geographic value of said geographical class, where a node class lower in said hierarchical data structure represents a geographical class that is hierarchically subordinate to and wholly encompassed in the geographical classes represented by node classes higher in said hierarchical data structure, and where said links between said node classes indicate which geographic values listed in a lower node class are geographically encompassed in a geographic value listed in a higher node class immediately adjacent to and above said lower node class in said hierarchical data structure;

identifying among the plurality of distinct <u>hierarchical data structures</u> one of said distinct <u>hierarchical data structures</u> having at least one geographical class in common, where a distinct <u>hierarchical data structure</u> is one which is hierarchically independent from all other <u>hierarchical data structures</u> and which has at least one geographic class that is unique to that hierarchical data structure;

merging portions of said distinct <u>hierarchical data structures</u> having one or more geographic classes in common by linking said structures to a single node class for each said geographical class said structures have in command; and

creating for each said node class in the merged hierarchical data structure a table

comprising an identifier describing the geographical class represented by the node class, an enumeration list encompassing the geographical values of said node class, and a plurality of linkage structures, where each linkage structure comprises an adjacent class identifier describing a geographical class immediately adjacent to and above said node class in said merged hierarchical data structure, and an adjacent link list, specifying for each first geographic value in said enumeration list, a second geographic value in an adjacent node class in which said first geographic value is geographically encompassed, and wherein the navigational field of a lowest geographical class includes said unique identifier thereby pointing to a corresponding one of said informational fields containing said same unique identifier.

2. The method of claim 1, wherein said step of ordering the navigational fields into a plurality of hiterarchical data structures further comprises the steps of:

creating a lowest node class for each <u>hierarchical data structure</u>, where said lowest node class encompasses nodes representing the unique identifiers; and

linking said lowest node class to the node class lowest in the <u>hierarchical data</u> structure which encompasses geographic values to indicate a relationship between said geographic values and the records of said informational fields.

3. A method for <u>converting</u> organized data from a data record structure into a hybrid structure which allows for efficient storage and retrieval of the organized data by a digital computer, where the data record structure includes a plurality of records each divided into a plurality of fields, where each field of the plurality of fields stores data of a field type associated with the field and where the organized data in the data record structure is encompassed in the plurality of fields of the plurality of records, the method comprising the steps of:

ordering said navigational fields into a plurality of distinct hierarchical data structures each comprising a plurality of node classes and links therebetween, where each said node class corresponds to a specific class of values and each node in said node class corresponds to a distinct value of said class, where a node class lower in said hierarchical data structure represents a class that is hierarchically subordinate to and wholly encompassed in the classes represented by node classes higher in said hierarchical data structure, and where said links between said node classes indicate which values listed in a lower node class are specifically related to a value listed in a higher node class immediately adjacent to and above said lower node class in said hierarchical data structure;

identifying among the plurality of distinct hierarchical data structures ones of said distinct hierarchical data structures having at least one class in common, where a distinct hierarchical data structure is one which is hierarchically independent from all other hierarchical data structures and which has at least one class that is unique to that hierarchical data structure;

merging portions of said distinct hierarchical data structures having one or more classes in common by linking said structures to a single node class for each class said structures have in common, where each said node in the merged node class represents the unique occurrence of a value encompassed in the common, yet unmerged node classes; and

creating for each said node class in the merged <u>hierarchical data structure</u> a table comprising an identifier describing the class represented by the node class, an enumeration list encompassing the values of said node class, and a plurality of linkage structures, where each linkage structure comprises an adjacent class identifier describing a class immediately adjacent to and above said node class in said merged <u>hierarchical data structure</u>, and an adjacent link list, specifying for each first value in said enumeration link, a second value in an adjacent node class to which said first value is related, and wherein the navigational field of a lowest class includes said unique identifier thereby pointing to a corresponding one of said informational fields containing said same unique identifier.

4. The method of claim 3, wherein said step of ordering the navigational fields into a plurality of <u>hierarchical data structures</u> further comprises the steps of:

creating a lowest node class for each <u>hierarchical data structure</u>, where said lowest node class encompasses nodes representing the unique identifiers; and

linking said lowest node class to the node class lowest in the hierarchical data

structure which encompasses a class of specific values to indicate a relationship between said specific values and the records of said informational fields.